

CS103  
WINTER 2025



# Lecture 03: **Propositional Logic**

## ***Logistics:*** Lecture Participation

# Lecture Participation

- Starting Wednesday, we will be using the website PollEV to ask questions in lecture for attendance credit.
- If you answer these questions in lecture, you'll get attendance credit for the day.
  - You don't need to have the right answers - you just need to respond to the questions.
- CGOE students: We automatically opt you out of participation, since we assume you aren't physically here.
- If you'd prefer not to attend lectures, that's okay! You can opt to count your final exam in place of participation.
  - We'll send out a form where you can opt-out of participation in Week 4.



Do not miss this deadline!

# Lecture Participation

- We'll dry-run PollEV questions today.
- Let's start with the following warm-up:

Make a music recommendation!

Answer at

<https://cs103.stanford.edu/pollev>

Click "Register" and enter your **Stanford e-mail** to get to the SUNet login page.

- Here are a few music recs of our own:
  - Jami Sieber - *Timeless*.
  - Aaron Parks - *Little Big* and *Little Big II*.
  - Arthur Moon - NPR Music Tiny Desk Concert.
  - Shakey Graves - *Roll the Bones* (check out *Audiotree Live* version).

Also:

pollev.com/103

# Propositional Logic

***Question:*** How do we formalize the definitions and reasoning we use in our proofs?

# Where We're Going

- ***Propositional Logic*** (Today)
  - Reasoning about Boolean values.
- ***First-Order Logic*** (Wednesday/Friday)
  - Reasoning about properties of multiple objects.

# Outline for Today

- ***Propositional Variables***
  - Booleans, math edition!
- ***Propositional Connectives***
  - Linking things together.
- ***Truth Tables***
  - Rigorously defining connectives.
- ***Simplifying Negations***
  - Mechanically computing negations.

# Propositional Logic

*TakeMath51*  $\vee$  *TakeCME100*

$\neg$ *FirstSucceed*  $\rightarrow$  *TryAgain*

*IsCardinal*  $\wedge$  *IsWhite*

*TakeMath51*  $\vee$  *TakeCME100*

$\neg$  *FirstSucceed*  $\rightarrow$  *TryAgain*

*IsCardinal*  $\wedge$  *IsWhite*

These are ***propositional variables***. Each propositional variable stands for a ***proposition***, something that is either true or false.

*TakeMath51*  $\vee$  *TakeCME100*

$\neg$ *FirstSucceed*  $\rightarrow$  *TryAgain*

*IsCardinal*  $\wedge$  *IsWhite*

These are ***propositional connectives***, which link propositions into larger propositions

# Propositional Variables

- In propositional logic, individual propositions are represented by ***propositional variables***.
- Each variable can take one of two values: true or false. You can think of them as **bool** values.

# Propositional Connectives

- There are seven propositional connectives, five of which will be familiar from programming.
- First, there's the logical “NOT” operation:

$\neg p$

- You'd read this out loud as “not  $p$ .”
- The fancy name for this operation is ***logical negation***.

# Truth Tables

- A ***truth table*** is a table showing the truth value of a propositional logic formula as a function of its inputs.
- Let's examine the truth tables for the connectives we're exploring today!

“I don’t love cupcakes.”

***LoveCupcakes*** : I love cupcakes.

$\neg$ ***LoveCupcakes***

# Propositional Variables

- In propositional logic, individual propositions are represented by ***propositional variables***.
- Each variable can take one of two values: true or false. You can think of them as **bool** values.
- In a move that contravenes programming style conventions, propositional variables are usually represented as lower-case letters, such as  $p$ ,  $q$ ,  $r$ ,  $s$ , etc.
  - That said, there's nothing stopping you from using multiletter names!

“I don’t love cupcakes.”

***LoveCupcakes*** : I love cupcakes.

$\neg$ ***LoveCupcakes***

# Propositional Connectives

- There are seven propositional connectives, five of which will be familiar from programming.
- Next, there's the logical “AND” operation:

$p \wedge q$

- You'd read this out loud as “ $p$  and  $q$ .”
- The fancy name for this operation is ***logical conjunction***.

“It’s cardinal and white.”

*IsCardinal* : It’s cardinal.

*IsWhite* : It’s white.

*IsCardinal*  $\wedge$  *IsWhite*

# Propositional Connectives

- There are seven propositional connectives, five of which will be familiar from programming.
- Then, there's the logical “OR” operation:

**$p \vee q$**

- You'd read this out loud as “ $p$  or  $q$ .”
- The fancy name for this operation is ***logical disjunction***. This is an *inclusive* or.

“You must take Math 51 or CME 100.”

***TakeMath51*** : You must take Math 51.

***TakeCME100*** : You must take CME 100.

***TakeMath51***  $\vee$  ***TakeCME100***

# Propositional Connectives

- There are seven propositional connectives, five of which will be familiar from programming.
- There's also the “truth” connective:

T

- You'd read this out loud as “true.”
- Although this is technically considered a connective, it “connects” zero things and behaves like a variable that's always true.

# Propositional Connectives

- There are seven propositional connectives, five of which will be familiar from programming.
- Finally, there's the “false” connective.

⊥

- You'd read this out loud as “false.”
- Like  $\top$ , this is technically a connective, but acts like a variable that's always false.

# Inclusive and Exclusive OR

- The  $\vee$  connective is an *inclusive* “or.” It's true if at least one of the operands is true.
  - It's similar to the `||` operator in C, C++, Java, etc. and the `or` operator in Python.
- Sometimes we need an *exclusive* “or,” which isn't true if both inputs are true.
- We can build this out of what we already have.

Write a propositional logic formula for the exclusive OR of  $p$  and  $q$ .

Answer at

<https://cs103.stanford.edu/pollev>

## *Quick Question:*

What would I have to show you to convince you that the statement  $p \wedge q$  is false?

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What would I have to show you to convince you that the statement  $p \vee q$  is false?

# de Morgan's Laws

$\neg(p \wedge q)$  *is equivalent to* ??

$\neg(p \vee q)$  *is equivalent to* ??

# de Morgan's Laws in Code

- **Pro tip:** Don't write this:

```
if (!(p() && q())) {  
    /* ... */  
}
```

- Write this instead:

```
if (!p() || !q()) {  
    /* ... */  
}
```

- (This even short-circuits correctly: if p() returns false, q() is never evaluated.)

# Mathematical Implication

# Implication

- We can represent implications using this connective:

$$p \rightarrow q$$

- You'd read this out loud as “ $p$  implies  $q$ .”
  - The fancy name for this is the ***material conditional***.
- ***Question:*** What should the truth table for  $p \rightarrow q$  look like?

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

An implication is false only when the antecedent is true and the consequent is false.

Every formula is either true or false, so these other entries have to be true.

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

***Important observation:***

The statement  $p \rightarrow q$  is true whenever  $p \wedge \neg q$  is false.

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

An implication with a false antecedent is called ***vacuously true***.

An implication with a true consequent is called ***trivially true***.

$p$	$q$	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T

***Please commit this table to memory.*** We're going to need it, extensively, over the next couple of weeks.

“If at first you don’t succeed, try again.”

***FirstSucceed*** : You succeed at first.

***TryAgain*** : You ought to try again.

$\neg \text{FirstSucceed} \rightarrow \text{TryAgain}$



IF IT AIN'T FRESHLY SLICED,  
IT AIN'T



OUTFRONT

2036



***JerseyMikes*** : It's Jersey Mike's.

***FreshlySliced*** : It's freshly sliced.

$\neg \text{FreshlySliced} \rightarrow \neg \text{JerseyMikes}$

Contrapositive?



***JerseyMikes*** : It's Jersey Mike's.

***FreshlySliced*** : It's freshly sliced.

$\neg \text{FreshlySliced} \rightarrow \neg \text{JerseyMikes}$

$\text{JerseyMikes} \rightarrow \text{FreshlySliced}$

# An Important Equivalence

- The truth table for  $p \rightarrow q$  is chosen so that the following is true:

$$p \rightarrow q \quad \text{is equivalent to} \quad \neg(p \wedge \neg q)$$

- Later on, this equivalence will be incredibly useful:

$$\neg(p \rightarrow q) \quad \text{is equivalent to} \quad p \wedge \neg q$$

## Side Note: Contrapositive

We can use truth tables to demonstrate the equivalence of  $p \rightarrow q$  and  $\neg q \rightarrow \neg p$ .

$p$	$q$	$p \rightarrow q$	$\neg q$	$\neg p$	$\neg q \rightarrow \neg p$
F	F	T	T	T	T
F	T	T	F	T	T
T	F	F	T	F	F
T	T	T	F	F	T

same :)



# The Biconditional Connective

# The Biconditional Connective

- In our previous lecture, we saw that the statement “ $p$  if and only if  $q$ ” means both that  $p \rightarrow q$  and  $q \rightarrow p$ .
- We can write this in propositional logic using the ***biconditional*** connective:

**$p \leftrightarrow q$**

- This connective’s truth table has the same meaning as “ $p$  implies  $q$  and  $q$  implies  $p$ .”
- Based on that, what should its truth table look like?

# Biconditionals

- The biconditional connective  $p \leftrightarrow q$  has the same truth table as  $(p \rightarrow q) \wedge (q \rightarrow p)$ .
- Here's what that looks like:

$p$	$q$	$p \leftrightarrow q$
F	F	T
F	T	F
T	F	F
T	T	T

One interpretation of  $\leftrightarrow$  is to think of it as equality: the two propositions must have equal truth values.

# Negating a Biconditional

- How do we simplify

$$\neg(p \leftrightarrow q)$$

using the tools we've seen so far?

- There are many options, but here are our two favorites:

$$p \leftrightarrow \neg q$$

$$\neg p \leftrightarrow q$$

Question to ponder: what is the truth table for these statements, and where have you seen it before?

# Operator Precedence

- How do we parse this statement?

$$\neg x \rightarrow y \vee z \rightarrow x \vee y \wedge z$$

- Operator precedence for propositional logic:

$\neg$   
 $\wedge$   
 $\vee$   
 $\rightarrow$   
 $\leftrightarrow$

- All operators are right-associative.
- We can use parentheses to disambiguate.

# Operator Precedence

- The main points to remember:
  - $\neg$  binds to whatever immediately follows it.
  - $\wedge$  and  $\vee$  bind more tightly than  $\rightarrow$ .
- We will commonly write expressions like  $p \wedge q \rightarrow r$  without adding parentheses.
- For more complex expressions, we'll try to add parentheses.
- Confused? ***Please ask!***

# The Big Table

Connective	Read Aloud As	C++ Version	Fancy Name	Negation
$\neg p$	“not”	!	Negation	$p$
$p \wedge q$	“and”	<code>&amp;&amp;</code>	Conjunction	$\neg p \vee \neg q$ $p \rightarrow \neg q$
$p \vee q$	“or”	<code>  </code>	Disjunction	$\neg p \wedge \neg q$
$\top$	“true”	<code>true</code>	Truth	$\perp$
$\perp$	“false”	<code>false</code>	Falsity	$\top$
$p \rightarrow q$	“implies”	<i>see PS2!</i>	Implication	$p \wedge \neg q$
$p \leftrightarrow q$	“if and only if”	<i>see PS2!</i>	Biconditional	$p \leftrightarrow \neg q$ $\neg p \leftrightarrow q$

Time-Out for Announcements!

# Submitting Work

- All assignments should be submitted through GradeScope.
  - The programming portion of the assignment is submitted separately from the written component.
  - The written component **must** be typed; handwritten solutions don't scan well and get mangled in GradeScope.
- All assignments are due at 1:00PM. You have three "late days" you can use throughout the quarter. Each automagically extends assignment deadlines from Friday at 1:00PM to Saturday at 1:00PM; at most one late day can be used per assignment.
  - **Very good idea:** Leave at least two hours buffer time for your first assignment submission, just in case something goes wrong.
  - **Very bad idea:** Wait until the last minute to submit.
- Your score on the problem sets is the square root of your raw score. So an 81% maps to a 90%, a 50% maps to a 71%, etc. This gives a huge boost even if you need to turn something in that isn't done.

# Office Hours

- Office hours have started (as of Sunday)! Think of them as “drop-in help hours” where you can ask questions on problem sets, lecture topics, etc.
  - Check the Guide to Office Hours on the course website for the schedule.
- TA office hours are held in person in the Huang basement. Sean’s are in Durand 331-B.
- Once you arrive, sign up through the CS Office Hours Queue so that we can help people in the order they arrived:

<https://queue.cs.stanford.edu/>

- Office hours are *much* less crowded earlier in the week than later. Stop by on Sunday, Monday, and Tuesday!

Back to CS103!

# Recap So Far

- A ***propositional variable*** is a variable that is either true or false.
- The ***propositional connectives*** are
  - Negation:  $\neg p$
  - Conjunction:  $p \wedge q$
  - Disjunction:  $p \vee q$
  - Truth:  $\top$
  - Falsity:  $\perp$
  - Implication:  $p \rightarrow q$
  - Biconditional:  $p \leftrightarrow q$

# Why All This Matters

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- Suppose we want to prove the following statement:  
“If  $x + y = 16$ , then  $x \geq 8$  or  $y \geq 8$ ”

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- Suppose we want to prove the following statement:

“If  $x + y = 16$ , then  $x \geq 8$  or  $y \geq 8$ ”

$$x < 8 \wedge y < 8 \rightarrow x + y \neq 16$$

“If  $x < 8$  and  $y < 8$ , then  $x + y \neq 16$ ”

**Theorem:** If  $x + y = 16$ , then  $x \geq 8$  or  $y \geq 8$ .

**Proof:** We will prove the contrapositive, namely, that if  $x < 8$  and  $y < 8$ , then  $x + y \neq 16$ .

Pick  $x$  and  $y$  where  $x < 8$  and  $y < 8$ . We want to show that  $x + y \neq 16$ . To see this, note that

$$\begin{aligned}x + y &< 8 + y \\&< 8 + 8 \\&= 16.\end{aligned}$$

This means that  $x + y < 16$ , so  $x + y \neq 16$ , which is what we needed to show. ■

# Why This Matters

- Propositional logic lets us symbolically manipulate statements and theorems.
  - This can help us better understand what a theorem says or what a definition means.
- It's also very useful for proofs by contradiction and contrapositive.
- Being able to negate statements mechanically can reduce the likelihood of taking an negation of contrapositive wrong.

# Negation Practice

- Here's a propositional formula that contains some negations. Simplify it as much as possible:

$$\neg(p \wedge q \rightarrow r \vee s)$$

# Negation Practice

- Here's a propositional formula that contains some negations. Simplify it as much as possible:

$$p \wedge q \wedge \neg r \wedge \neg s$$

# Negation Practice

- Here's a propositional formula that contains some negations. Simplify it as much as possible:

$$\neg((p \vee (q \wedge r)) \leftrightarrow (a \wedge b \wedge c \rightarrow d))$$

# Negation Practice

- Here's a propositional formula that contains some negations. Simplify it as much as possible:

$$(p \vee (q \wedge r)) \leftrightarrow (a \wedge b \wedge c \wedge \neg d)$$

# Next Time

- *First-Order Logic*
  - Reasoning about groups of objects.
- *First-Order Translations*
  - Expressing yourself in symbolic math!